

# Queue

Data Structure

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# Introduction

- It is linear data structure
- It is collection of items – List
- Queue means line of items waiting for their turn
- Queue has two ends
  - Elements are added at one end.
  - Elements are removed from the other end.

# Queue

- Removal of data item is restricted at one end known as FRONT
- Insertion of data item is restricted at other end known as REAR
- The FRONT and REAR are used in describing a linear list, when queue is implemented
- First element inserted in list, will be the first to be removed - FIFO

# Queue as FIFO

- The element inserted first will be removed first from the Queue
- Thus, Queue is known as FIFO (**F**irst **I**n-**F**irst **O**ut) or FCFS (**F**irst **C**ome **F**irst **S**erve)
- **Examples of Queue**
  - People waiting in Queue to purchase tickets at railway station or cinema hall, where the first person in the queue will be served first

# Representation of Queue

- It has two pointer variables
  - FRONT : Containing the location of the front element of the queue
  - REAR : Containing the location of the rear element of the queue
- When queue is empty  
FRONT = -1 and REAR = -1

# Operations of Queue

- Insertion

- Adding an element in queue will increased value of REAR by 1

- $REAR = REAR + 1$

- Removal

- Removing an element from queue will increased value of FRONT by 1

- $FRONT = FRONT + 1$

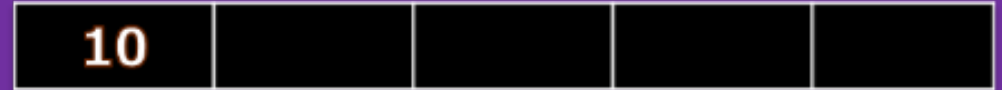
**Queue  
Empty**

FRONT = -1, REAR =  
-1



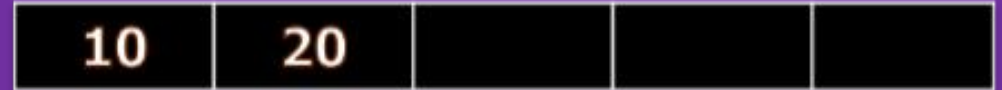
ADD(Q,10)

FRONT = 0 ,REAR =  
0



ADD(Q,20)

FRONT = 0 ,REAR  
= 1



ADD(Q,30)

FRONT = 0 ,REAR  
= 2



ADD(Q,40)

FRONT = 0 ,REAR  
= 3



ADD(Q,50)

FRONT = 0 ,REAR  
= 4



Remove(Q)

FRONT = 1 ,REAR  
= 4



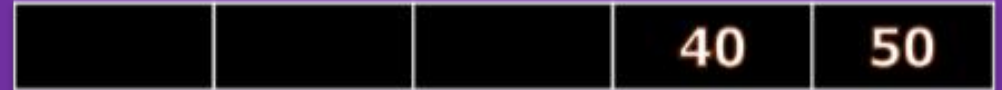
Remove(Q)

FRONT = 2 ,REAR  
= 4



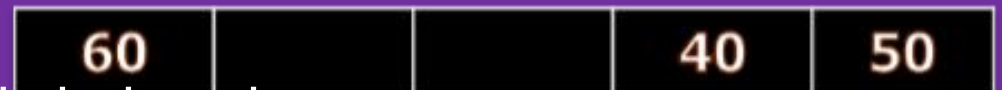
Remove(Q)

FRONT = 3 ,REAR  
= 4



ADD(Q,  
60)

FRONT = 3 ,REAR  
= 1



# Types of Queue

- Queue as Array
- Circular Queue
- Priority Queue
- Input Restricted Queue
- Output Restricted Queue
- Dqueue



## Queue as Array : Insert

Initially when the QUEUE is empty, set FRONT = NULL and REAR = 0

Step 1: start

Step2: [check for queue is over flow or not]

If (REAR >n) or (REAR==FRONT)

Print “queue is overflow”

else

go to step 3

Step 3: [enter the item]

QUEUE[REAR]=value

REAR=REAR+1

Step 4:[ check condition]

If(FRONT==null)

FRONT=0

Step 5:end

## Queue as Array : Delete

```
Step 1: start
Step 2: [check for queue is under flow or not]
If front>N or front==Null
Print "queue is underflow"
else
goto step 3
Step 3: [check condition]
If front==rear
Front==null
Rear=0
else
goto step 4
Step 4: [delete element]
Queue[front]=null
Step 5: front=front+1
Step 6: end
```

# Circular Queue

- To solve this problem, queues implement wrapping around. Such queues are called Circular Queues.
- Both the front and the rear pointers wrap around to the beginning of the array.
- It is also called as “Ring buffer”.
- Items can inserted and deleted from a queue in  $O(1)$  time.

# Circular Queue

- When a new item is inserted at the rear, the pointer to rear moves upwards.
- Similarly, when an item is deleted from the queue the front arrow moves downwards.
- After a few insert and delete operations the rear might reach the end of the queue and no more items can be inserted although the items from the front of the queue have been deleted and there is space in the queue.

# QINSERT(Queue, N, FRONT, REAR, ITEM)

1 [Queue already filled ?]

if  $FRONT = 0$  and  $REAR = N-1$ , or if  $FRONT=REAR + 1$  then  
write : Overflow and Return

2 [Find new value of REAR]

if  $FRONT=-1$  then [Queue initially empty]

Set  $FRONT = 0$  and  $REAR = 0$

else if  $REAR = N-1$  then

Set  $REAR = 1$

else

Set  $REAR = REAR + 1$

[End of If structure]

3 Set  $QUEUE[REAR] = ITEM$  [This inserts new element]

4 Return

# QDELETE(Queue, N, FRONT, REAR, ITEM)

1 [Queue already empty]

If FRONT = -1, then Write : Underflow and Return

2 Set ITEM = Queue[FRONT]

3 [Find new value of FRONT]

If FRONT = REAR, then [Queue has only one element to start]

Set FRONT = -1 and REAR = -1

Else if FRONT = N-1, then

Set FRONT = 0

Else

Set FRONT = FRONT + 1

[End of If Structure]

4 Return

# Priority Queue

Suppose that you have a few assignments from different courses. Which assignment will you want to work on first?

| Course                       | Priority | Due day      |
|------------------------------|----------|--------------|
| Database Systems             | 2        | October 3    |
| CONMS                        | 4        | October 10   |
| Advance 'C' & Data Structure | 1        | September 29 |
| Software Engineering         | 3        | October 7    |
| C++                          | 5        | October 15   |

You set your priority based on due days. Due days are called keys.

# Priority Queue

- It is collection of elements where elements are stored according to the their priority levels
- Inserting and removing of elements from queue is decided by the priority of the elements
- The two fundamental methods of a priority queue  $P$ :
  - insertItem( $k, e$ ): Insert an element  $e$  with key  $k$  into  $P$ .
  - removeMin(): Return and remove from  $P$  an element with the smallest key.



# Priority Queue

- When you want to determine the priority for your assignments, you need a value for each assignment, that you can compare with each other.
- **key**: An object that is assigned to an element as a specific attribute for that element, which can be used to identify, rank, or weight that element.

# Example: Student records

Any of the attributes Student Name, Student Number, or Final Score can be used as the primary key.

Note: Ke

| Student Name | ID       | Final Score (out of 450) |
|--------------|----------|--------------------------|
| Ashwin       | 09BCA08  | 310                      |
| Payal        | 09BCA80  | 311                      |
| Darshika     | 09BCA24  | 380                      |
| Nilkamal     | 09BCA75  | 400                      |
| Nikunj       | 09BCA74  | 440                      |
| Mori         | 09BCA102 | 400                      |

# Rules to maintain a Priority Queue

- The elements with the higher priority will be processed before any element of lower priority
- If there are elements with the same priority, then the element added first in the queue would get processed

# Priority Queue Insert

PQInsert (M, Item)

Step 1 Find the Row M

Step2 [Reset the Rear Pointer]

If  $\text{Rear}[M] = N-1$  then  $\text{Rear}[M] = 0$

Else

$\text{Rear}[M] = \text{Rear}[M]+1$

Step 3 [Overflow]

If  $\text{Front}[M] = \text{Rear}[M]$  then Write (“This Priority Queue is full”)

Return

Step 4 [Insert Element]

$Q[M][\text{Rear}[M]] = \text{Item}$

Step 5 [Is Front Pointer Properly Set]

If  $\text{Front}[M] = -1$  then  $\text{Front}[M] = 0$

Return

Step 6 Exit

# Priority Queue Delete

PQDelete (K, Item)

Step 1 Initialize  $K = 0$

Step 2 while ( $\text{Front}[K] = -1$ )

$K = K+1$

[To find the first non empty queue]

Step 3 [Delete Element]

Item =  $Q[K][\text{Front}[K]$

Step 4 [Queue Empty]

If  $\text{Front}[K] = N-1$  then  $\text{Front}[K] = 0$

Else

$\text{Front}[K] = \text{Front}[K]+1$

Return Item

Step 6 Exit

# Deque

- Deque stands for double-end queue
- A data structure in which elements can be added or deleted at either the front or rear
- But no changes can be made in the list
- Deque is generalization of both stack and queue

# Removing Element from Deque

- There are two variations of a deque. These are
  - Input Restricted Deque
    - An input restricted deque restricts the insertion of elements at one end only, but the deletion of elements can perform at both the ends.
  - Output Restricted Deque
    - An output restricted queue, restricts the deletion of elements at one end only, and allows insertion to be done at both the ends of deque

# Possibilities

- The two possibilities that must be considered while inserting or deleting elements into the queue are :
  - When an attempt is made to insert an element into a deque which is already full, an **overflow occurs**.
  - When an attempt is made to delete an element from a deque which is empty, **underflow occurs**.



# Representation of Deque



# Inserting Element in Deque

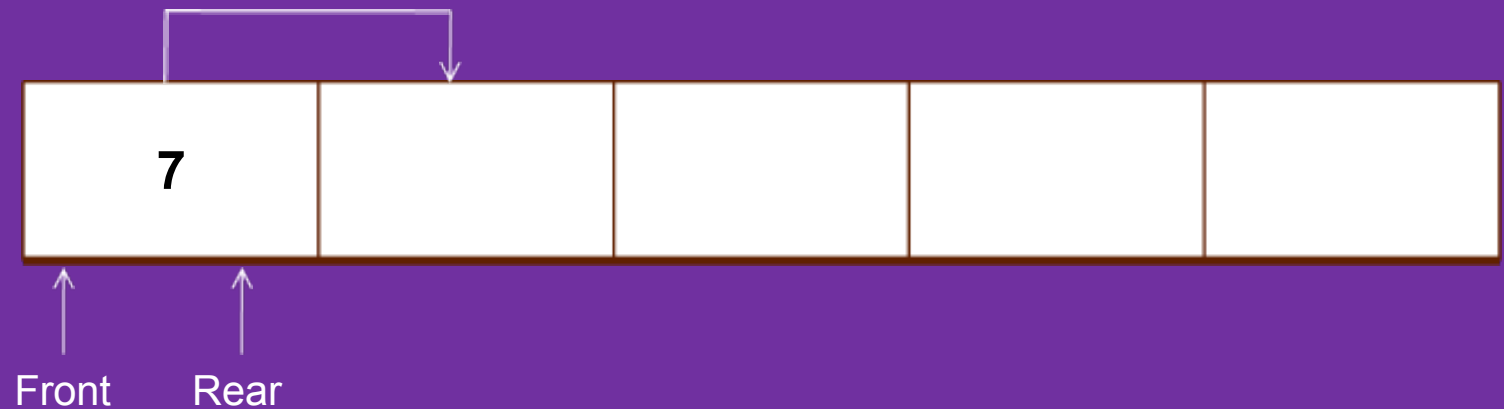
InsertQatEnd(dq,7)



InsertQatBeg(dq,10)



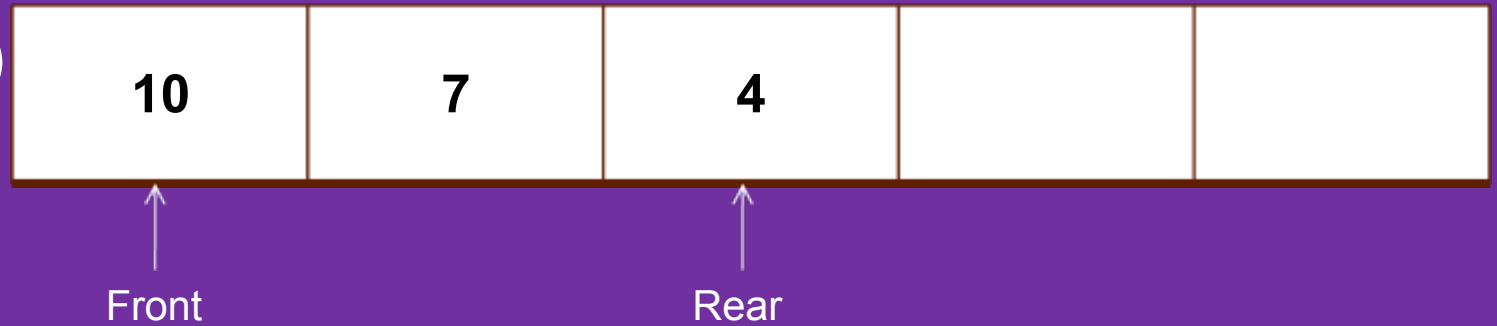
InsertQatBeg(dq,  
10)



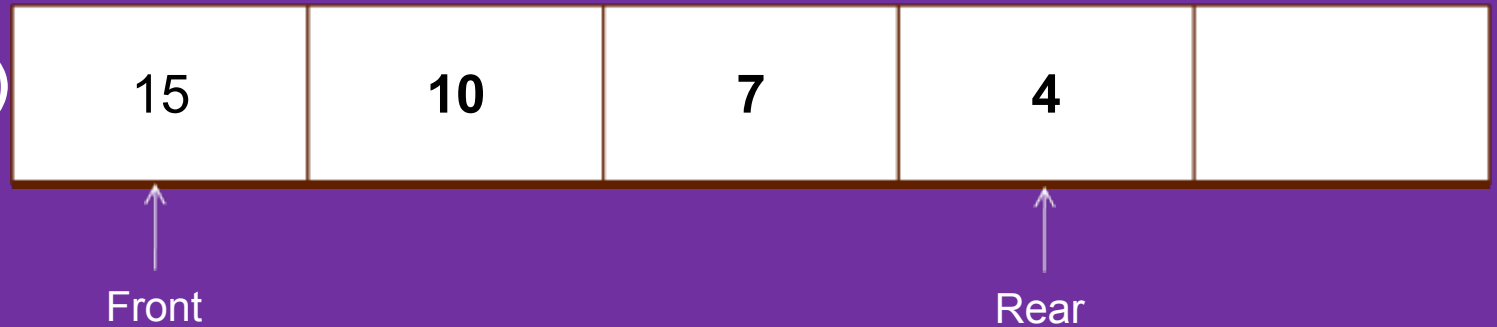
**InsertQatBeg(dq,10)**



**InsertQatEnd(dq,4)**



**InsertQatBeg(dq,15)**

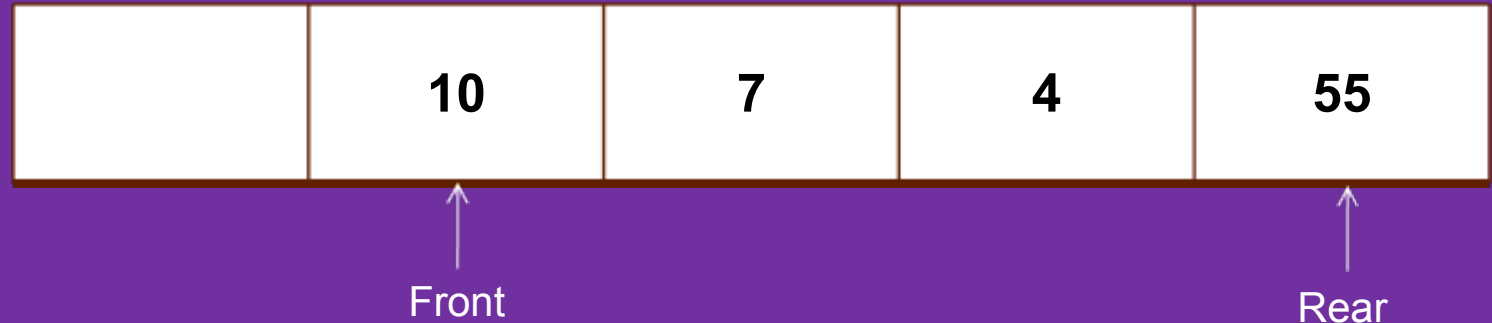


**InsertQatEnd(dq,55)**

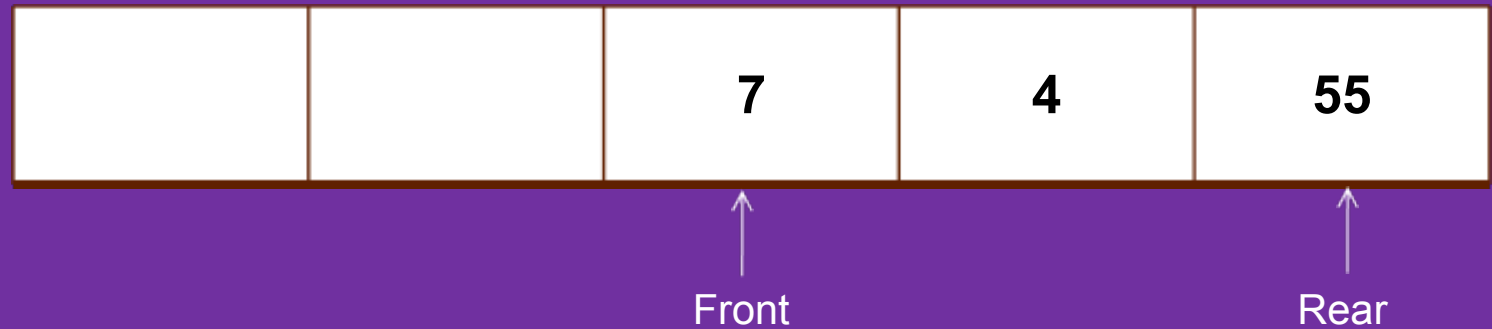


# Removing Element from Deque

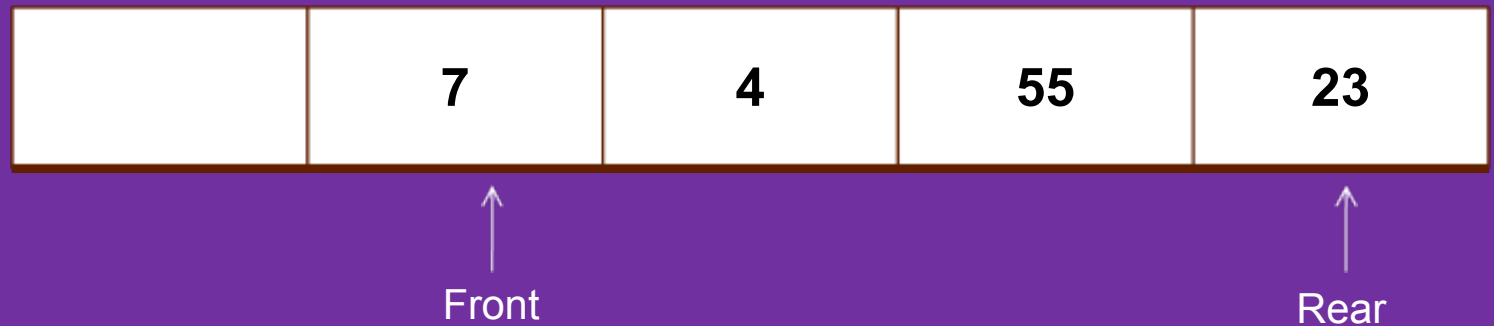
delQBeg()



delQBeg()



InsertQatEnd(dq,23)



# Deque : Insert

- There are two variations of a deque. These are
  - Input Restricted Deque
    - An input restricted deque restricts the insertion of elements at one end only, but the deletion of elements can perform at both the ends.
  - Output Restricted Deque
    - An output restricted queue, restricts the deletion of elements at one end only, and allows insertion to be done at both the ends of deque

# Input Restricted queue : Insert Element

Step 1: start

Step 2:[check condition for overflow]

If(rear==N-1 && front==0 or front=rear+1)

Print "over flow"

else

goto step 3

Step 3: [check condition]

If(front==null)

front = 0 and rear=-1

else

goto step 4

Step 4: [check condition and value]

If (rear== N -1)

rear=0

Else

rear=rear+1

Step 5: [Add value]

dq[rear]=value

Step 6:end

# Input Restricted queue : Delete Beginning

Step 1: start

Step 2: [check condition for underflow]

If(rear==null & front==null)

Print"underflow"

else

goto step 3

Step 3: [delete]

dq[rear]=null

Step 4: [check condition]

If(rear==front)

front=rear=null

else

rear--;

Step 5: end

# Input Restricted queue : Delete End

Step 1: start

Step2 : [check condition for under flow)

If(rear==null & front==null)

Print “under flow”

else

goto step 3

Step 3: [delete]

dq[front]=null

Step 4: [check condition]

If(rear==front)

rear=-1 and front=null

else

front=front+1

Step 5: end